

Multiparty Privacy Management and Cluster Analysis in Social media (MPMC)

Sreekutty J.S., Hamil Stanly, Archa A.T.

Abstract— The current mainstream Social Media infrastructures without multi-party privacy management do not preserve the privacy of each users as they are unable to appropriately control to whom co-owned items are shared. For solving this issue, here introduce a mediator which detects conflict and propose the appropriate solution. The solution based on some principle which checks whether the uploaded item may cause any detrimental action to one of the user involved in that. This mechanism inspecting all users' individual privacy policies and generate a comparison between them; if it detect at least two contradictory decision in the privacies, apply concession rule , instantiates based on users willingness to change an action(granting/denying).

Index Terms— Conflict, Multiparty privacy, Online social network, Social networking services.

I. INTRODUCTION

Social media are computer-mediated technologies that allow creating and sharing of information, ideas, career interests and other forms of expression via virtual communities and networks. Photos are a proven way to increase your engagement on social media. In terms of both likes and comments photos far exceeded other post types in engaging consumer responses. Only the uploader's privacy preferences are considered by the social media ;it does not concern about other user affected by that item. The affected users are the co-owners and they have their own privacy settings that sometimes it conflict uploaders privacy. The existing approaches needs manual negotiation techniques for solve this problem but it is very time consuming one.

Recent survey explains a number of important design considerations for photo privacy tools around the importance of identity and impression management and the tensions of ownership. "Restrict Others Tool" explicitly dealt with the natural tension that arises between the owner of the photo, and those tagged in it. Here created a lightweight means for users to negotiate desired sharing, complementing the existing privacy coping mechanisms that users currently employ. In manipulating these ownership tensions, and believe their tool would help users achieve more desired privacy while still maximizing the social value of sharing[2]. But this design technique only inspects 2 parties privacy policy but we need

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to set the members manually. This is the reason for an automatic computational mechanism that does not need any user interaction, so here suggests first computational mechanism which automatically inspects all the privacy preferences , if any conflict is detected, solve it.

II. OBJECTIVE

The main objective is to provide individual privacy policy settings to all users. Also provide privacy conflict detection and automated privacy conflict resolution

III. SYSTEM OVERVIEW

Number of individual using the social media increases dramatically . This system is to be set up for two types of users either having same privacy policies or having that of different. Firstly identifying the users and co-owner of an item. One of them wants to upload an item and other users become affected users for that same item. So here use a mediator which could be integrated as backend, it detect all the privacy policies of negotiating users and if any conflict found ,it flags off all. The proposed system includes following phases:

- User Activities
- Conflict Detection
- Conflict Resolution

1.USER ACTIVITIES

i. Registration

User can register into these media by entering his Name, DOB, Address, Phone num, Email Id, Password, Security questions etc. The registered user can login into the system by using this E-mail Id and Password.

ii. Data Uploading

Every user can upload the data especially pictures and comments that mention multiple users through these social networking services. User can provide an optional message or heading with this image. Image can be selected from the existing resources or taken instantly. Users wants more likes and share, Because they want to popularize themselves.

iii. Privacy settings

Select friends and add them to a group. Both uploader and users who all are affected by the item have their own privacy settings for it. They can decide to grant/deny access to the targeted user for that item.

Format, P = ({A} . {E})

Where, A represent the groups that are granted to access an item, E is the excluding member in the group that need policies.

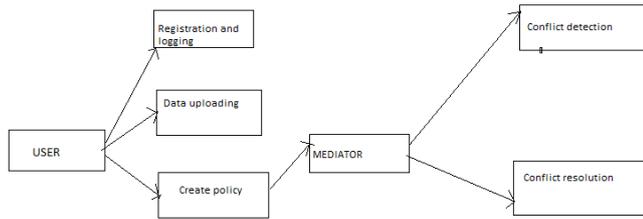


Fig.1: Mechanism overview

2. CONFLICT DETECTION

Mediator inspects all the privacy policies of each negotiating users. For comparing each privacy policies, consider the effect in which particular policies has on target user .The owner has availing two actions granted the access and deny the access. so here consider an action vector $v(t)$. if all the negotiating users assign same action to the particular item , it means that there is no conflict if else there occurs a conflict in the implementation section mediator runs the conflict detection algorithm and find out the conflict[1].

Algorithm 1 Conflict Detection

```

Input:  $N, P_{n_1}, \dots, P_{n_{|N|}}, T$ 
Output:  $C$ 
1: for all  $n \in N$  do
2:   for all  $t \in T$  do
3:      $v_n[t] \leftarrow 0$ 
4:     for all  $G \in P_n.A$  do
5:       if  $\exists u \in G, u = t$  then
6:          $v_n[t] \leftarrow 1$ 
7:       end if
8:     end for
9:   end for
10:  for all  $e \in P_n.E$  do
11:     $v_n[e] \leftarrow \neg v_n[e]$ 
12:  end for
13: end for
14:  $C \leftarrow \emptyset$ 
15: for all  $t \in T$  do
16:   Take  $a \in N$ 
17:   for all  $b \in N \setminus \{a\}$  do
18:     if  $v_a[t] \neq v_b[t]$  then
19:        $C \leftarrow C \cup \{t\}$ 
20:     end if
21:   end for
22: end for
    
```

For example, assume that Anju is the owner of a comment or picture uploaded (content), The content may contains other peoples such as Appu, Arya, Arun, Veena and Anusree. So Arya is one of the negotiating user in this content. Anju wants to share this content to her “Myfriends” group. It contains {Appu, Arun, Anusree}. But arya wants to share this content to her “Family”. It contains {Appu, Arun and Veena}.

Hence $P1 = \{ 1,0,1,0,1\}$
 $P2 = \{1,1,1,1,0\}$

This means that $P1 \neq P2$ (A conflict in both privacy policies).

3. CONFLICT RESOLUTION

The mediator suggest the solution based on three principles:

1. If an item may cause any detrimental action to one of the user involved , that item should not be shared

2. An item should be shared if it is not generating any harm to all negotiating users and also it is vital for one person
3. In the conflicting section, the mediator automatically solve the conflict based on concession rule.[1]

If the conflict is to be detected the mediator runs the conflict resolution algorithm and the output is generated.

Algorithm 2 Conflict Resolution

```

Input:  $N, P_{n_1}, \dots, P_{n_{|N|}}, C$ 
Output:  $\bar{o}$ 
1: for all  $c \in C$  do
2:
3:   if  $\forall n \in N, W(n,c)$  is HIGH then
4:      $o[c] \leftarrow \text{modified\_majority}(P_{n_1}, \dots, P_{n_{|N|}}, c)$ 
5:     continue
6:   end if
7:
8:   if  $\exists a \in N, W(a,c)$  is LOW then
9:     if  $\exists b \in N, W(b,c)$  is LOW  $\wedge v_a[c] \neq v_b[c]$  then
10:       $o[c] \leftarrow 0$ 
11:    else
12:       $o[c] \leftarrow v_a[c]$ 
13:    end if
14:  end if
15: end for
    
```

For solving the conflict, the mediator first inspect all negotiating users privacy preferences and determine what makes the users to change their decision(grant or deny the access). For this purpose, the mediator first inspect the sensitivity of an item and Prioritization of an item.

a. Sensitivity of an item

If an item is very sensitive for a user, he/she will not share this item to others. For example, Ammu is a co-owner of a photo sharing activity. Her friend Achu Upload a picture regarding a night party which shows the intimate relationship between Achu and Ammu or comment that mention about Ammu. Achu and Ammu are cousins. Achu set the privacy preference as his “Family”. It contains Ammus Mother. Ammu doesn’t want to disclose this secret to her mother. So, for Ammu, this item is very sensitive to her.

The sensitivity of an item is to be calculated by

$$S_n = (1/G_n) \sum T_n(G)[1]$$

Where S_n = Sensitivity of an item

$T_n(G)$ = Tie strength between two parties

G_n = Groups

Tie strength between groups can be determined by estimating the minimum strength needed for an item to shared in a group $\rho(n,c)$.

$$S_{\text{ammu}} = (1/G_{\text{ammu}}) \sum T_{\text{ammu}}(\text{family})$$

$$T_{\text{ammu}}(\text{family})=5.$$

b. Prioritization of an item

This section deals with conflicting areas. It determines about different negotiating users assigns different actions to the same target user. For example, Appu planned to provide a surprise party to his close friend Vinu. He wants to invite all his other close friends to Vinu home except Vinu. The relationship between the user and target user is high. Geethu is another member in this group. She is Vinu's sister and she

disclose every matter to her brother, she doesn't want to hide anything from her brother. The relationship between Geethu and Vinu is high. Hence prioritization can be determined by the formula,

$$P_n^{\odot} = S_n - \dot{\rho}(n,c) \quad [1]$$

Where P_n^{\odot} = Prioritization of an item

S_n = Sensitivity of an item

$\dot{\rho}(n,c)$ = Minimum strength needed for an item to shared in a group

$\dot{\rho}(n,c)$ can be calculated by an example:

where n denotes negotiating user and c denotes target user.

If Anju wants to share an item to three peoples – Anu, Achu, Allu. Anu is her best friend, Achu is her good friend and allu is her friend.so

$$\dot{\rho}(\text{Anju, Anu}) = 5$$

$$\dot{\rho}(\text{Anju, Achu}) = 4$$

$$\dot{\rho}(\text{Anju, Allu}) = 3.$$

c. Determine the Desire

The desire to change an action can be determined by above two criteria.

If sensitivity of an item and its prioritization is very high-the negotiating user wont share the item to target user. O the contradictory, if sensitivity of an item and its prioritization is very loe-the negotiating user will ready to share the item to the target user.

$$D = 0.5 \left(\frac{(\partial - P_n^{\odot})}{(\partial + P_n^{\odot})} \right) / \left(\frac{(\partial - S_n)}{(\partial + S_n)} \right)$$

Where D = Determine the Desire

∂ = Tie strength maximum value.

In the above example tie strength maximum value goes to 5(Best friends).

4. CLUSTER ANALYSIS

Cluster analysis is used in this media for grouping all the photos that uploaded in such a way that these pictures are in the same group(called a cluster) are more similar(in sense of their colour, edge etc) to each other than to those in other groups(clusters). Here clustering is used to recognize communities within large groups of people. Centroid based clustering is used in which clusters are represented by a central vector, which may not necessarily be a member of the dataset. When all these pictures are considered to k, kmeans clustering gives optimization as: k cluster on center and assign the object to the nearest cluster center, such that the squared distance from the clusters are minimized.

Algorithmic steps for kmeans clustering

Let $X = \{x_1, x_2, \dots, x_n\}$ be the set of datapoints and $V = \{v_1, v_2, \dots, v_c\}$ be the set of centers.

- i. Randomly select "c" cluster centers.
- ii. Calculate the distance between each datapoint and cluster centers.
- iii. Assign the datapoint to the center whose distance from the cluster center is minimum of all the cluster centers.
- iv. Recalculate the new cluster using datapoints.

$$V_i = (1/c_i) \sum x_i$$

Where c_i represents the number of datapoints in ith cluster.

- v. Recalculate the distance between each datapoint and new obtained cluster centers.
- vi. If no datapoint was reassigned, then stop, otherwise repeat from step iii.

IV. CONCLUSION

In recent years social media got huge growth with billions of users but they also have to face privacy violations. Users uploads no of items on a social media and granting/denying access to other users for that particular item is based on their privacy preferences. Here system concerns uploader's as well as affected user's policy for that item to grant/deny access to the targeted users. For that, system provides new mechanism of conflicts identification with atomic solution on a social media. The system having a graphical representation which shown that the percentage of people who have contradictory decision in an area - who accept the solution and reject the solution.

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